The Division Algorithm for \mathbb{N} and \mathbb{Z}

Theorem (Division Algorithm for \mathbb{N}) Suppose *a* and *b* are natural numbers and that $b \leq a$. Then there is a natural number *q* and a whole number *r* such that a = bq + r and $0 \leq r < b$. Moreover, *q* and *r* are <u>unique</u>.

(We usually call q the "quotient" and r the "remainder" when a is divided by b.)

Proof Let $A = \{s \in \mathbb{N} : sb > a\}$. There must be at least one such value of *s* (*see the* "*Archimedean Principle*", *text*, *p*. 105), so $A \neq \emptyset$. By the Well-Ordering Principle (WOP), *A* contains a <u>smallest</u> element: call it *l*.

Since $1 \notin A$, we know that l > 1 so q = l - 1 is a natural number. Define r = a - bq. This makes a = bq + r. (But we still need to prove that $0 \le r < b$).

We know that $\begin{cases} lb > a & \text{because } l \in A \\ (l-1)b = qb \le a & \text{because } q < l \text{ and therefore } q \notin A. \end{cases}$

Since $qb \leq a$, we have $r = a - bq \geq 0$.

To see that r < b: if $r = a - bq \ge b$, we would have $a \ge b + bq = b + b(l - 1) = lb$, which is false.

Therefore $0 \le r < b$ so the proof "there exist" q and r with the desired properties is complete

<u>To prove q and r are unique</u>: Suppose a = bq + r where $0 \le r < b$ (*) and that a = bq' + r' where $0 \le r' < b$ (**)

Subtracting these equations and rearranging gives b(q - q') = r' - r, and therefore

$$b|(q-q')| = |r'-r|.$$
 (***)

Also, we know from (*) that $-b < -r \le 0$, and adding this inequality to the inequality $0 \le r' < b$ (**) gives -b < r' - r < b so that

|r' - r| < b

Substituting this in (***), we get b|(q - q')| < b.

This means that |q - q'| < 1 and, since |q - q'| is a integer, this implies that |q - q'| = 0, that is, q = q'. From this, and the equations

$$a = bq + r$$
$$a = bq' + r'$$

we get that r = r'. Therefore q, r are unique.

(over \rightarrow)

We can also state a division algorithm for \mathbb{Z} .

Theorem (Division Algorithm for \mathbb{Z}) Suppose *a* and *b* are integers and that b > 0Then there is an integer *q* and an integer *r* such that a = bq + r and $0 \le r < b$. Moreover, *q* and *r* are unique.

There are various ways to state the division algorithm for \mathbb{Z} . In this version, we require that the divisor b > 0, so actually $b \in \mathbb{N}$, and that $0 \le r < b$. Another version (you might try to prove it) allows b to be any nonzero integer and has $0 \le r < |b|$. In all versions, the statement requires that the remainder r be nonnegative: that fact is usually what's important when the Division Algorithm is used.

Proof If a > 0, we get q and r from the Division Algorithm in \mathbb{N} .

If a = 0, let q = r = 0.

If a < 0, then apply the Division Algorithm in \mathbb{N} for dividing -a by b. There are natural numbers q' and r' for which

-a = bq' + r' where $0 \le r' < b$

Then a = b(-q') - r' where $-b < -r' \le 0$.

Using this equation:

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Case i) If -r' = 0:
Let q = -q', r = 0. Then a = b(-q') - r' = bq + r where 0 \le r < b
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Case ii) If -b < -r' < 0:
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Let q = -q' - 1 and r = b - r'.

Then
$$a = b(-q') - r' = b(-q') - b - r' + b$$

= $b(-q'-1) + (b - r')$
= $bq + r$, where $0 < b - r' = r < b$
(since $-b < -r' < 0$)

So, in both cases, we can find integers q and r for which a = bq + r, with 0 < r < b.

The proof that q and r are unique is left as an exercise (see proof of the previous theorem for *ideas*). •

Example The division algorithm in \mathbb{N} : 3 < 7 so we can write 7 = 3q + r where $0 \le r < 2$ (namely, with q = 2 and r = 1)

The division algorithm in \mathbb{Z} (in the form stated above, requiring the divisor b > 0) with b = 3 and a = -7 says that we can write -7 = 3q + r, where $0 \le r < 3$. Here the values that work are q = -3 and r = 2, and that's the <u>only</u> way to pick qand r if you want $0 \le r < 3$. We proved the Division Algorithm for \mathbb{N} using WOP. Here's an alternate proof using PMI, doing "induction on *b*." Stated more formally, we want to prove:

$$(\forall b \in \mathbb{N})(\forall a \in \mathbb{N})(\exists q \in \mathbb{N})(\exists r \in \mathbb{N}) \ (a = bq + r) \land (0 \le r < b)$$

To prove this universal statement (working left to right), we pick any $b \in \mathbb{N}$.

For this arbitrary but fixed natural number b we need to prove that

$$(\forall a \in \mathbb{N}) (\exists q \in \mathbb{N}) (\exists r \in \mathbb{N}) \ (a = bq + r) \land (0 \le r < b)$$

This is a statement of the form $(\forall a \in N) P(a)$, where P(a) is the statement

$$(\exists q \in \mathbb{N})(\exists r \in \mathbb{N}) \ (b = aq + r) \land (0 \le r < a)$$

We use induction on the natural number a.

This looks just a little odd, but a <u>is</u> a natural number, so induction is OK; if it makes you more comfortable, change "a" everywhere to "n."

Proof <u>Base case</u>: Suppose a = 1.

If
$$b = 1$$
, let $q = 1$ and $r = 0$. Then $a = bq + r$ and $0 \le r < 1$
If $b > 1$, let $q = 0$ and $r = 1$. Then $a = bq + r$ and $0 \le r < 1$.

So P(1) is true.

<u>Induction step</u>: Suppose P(a) is true for some particular value of a. Thus, we are assuming (for this value of a) that there are natural numbers q' and r' for which a = bq' + r' and $0 \le r' < b$.

We need to prove that P(a + 1) is true, that is, that there exists natural numbers q and r for which a + 1 = bq + r, where $0 \le r < b$.

<u>Case i</u>: If r' = b - 1: a = bq' + r' = bq' + b - 1, so a + 1 = bq' + b = b(q' + 1) + 0

> Let q = q' + 1 and r = 0Then a + 1 = bq + r, where $0 \le r < b$

<u>Case ii</u>: If r' < b - 1: a = bq' + r' so a + 1 = ab + (r' + 1)

Let
$$q = q'$$
 and $r = r' + 1$
Then $a + 1 = bq + r$, where $0 \le r < b$.

In both cases, we can find the necessary q and r. So P(k+1) is true.

By PMI, $(\forall a \in \mathbb{N}) P(a)$ is true (for the particular *b* we chose). Since the argument works no matter which $b \in \mathbb{N}$ we choose, we conclude that

$$(\forall b \in \mathbb{N})(\forall a \in \mathbb{N})(\exists q \in \mathbb{N})(\exists r \in \mathbb{N}) \ (a = bq + r) \land (0 \le r < b) \bullet$$