

Lecture 25: Bounding the remainder

Given a function f with continuous $(n+1)^{\text{st}}$ derivative in a neighborhood $N(a) = (a-c, a+c)$ of a , we derived the formula

$$E_n(x) := f(x) - T_n(x) = f(x) - \sum_{k=0}^n \frac{f^{(k)}(a)}{k!} (x-a)^k$$
$$\stackrel{\textcircled{=}}{=} \int_a^x \frac{f^{(n+1)}(t)}{n!} (x-t)^n dt$$

for the error in the n^{th} Taylor approximation (polynomial) at a .

This leads of once to the

or $[x, a]$

Theorem 1: $|E_n(x)| \leq \frac{\sup_{t \in [a, x]} |f^{(n+1)}(t)|}{(n+1)!} |x-a|^{n+1}$.

Proof: In general, we can bound integrals of continuous functions by reasoning that $-|F(t)| \leq F(t) \leq |F(t)| \Rightarrow -\int_a^b |F| dt \leq \int_a^b F dt \leq \int_a^b |F| dt$
 $\Rightarrow \left| \int_a^b F(t) dt \right| \leq \int_a^b |F(t)| dt$. Writing $M := \sup_{t \in [a, x]} |f^{(n+1)}(t)|$, we have

$$|E_n(x)| \leq \int_a^x \frac{|f^{(n+1)}(t)|}{n!} |x-t|^n dt \leq \int_a^x \frac{M}{n!} |x-t|^n dt = \frac{M}{(n+1)!} |x-a|^{n+1}. \quad \square$$

Corollary: Writing $M := \sup_{x \in N(a)} |f^{(n+1)}(x)|$, the error function $|E_n(x)| \leq \frac{M}{(n+1)!} |x-a|^{n+1}$ on $N(a)$. In particular, $\lim_{x \rightarrow a} \frac{E_n(x)}{(x-a)^n} = 0$.

Definition: If $g(x) \neq 0$ for $x \in N^*(a)$ and $\lim_{x \rightarrow a} \frac{f(x)}{g(x)} = 0$, then we write $f(x) = o(g(x))$.

So in the Corollary we have $E_n(x) = o(|x-a|^n)$.

Why might we care about bounding the error?

① So we can decide how closely $T_n(x)$ approximates $f(x)$ at particular values of x . For instance, with $f(x) = e^x$ & $a=0$,

$$\left| e - \sum_{k=0}^n \frac{1}{k!} \right| = |E_n(1)| \leq \frac{\sup_{t \in [0,1]} e^t}{(n+1)!} |1-0|^{n+1} = \frac{e}{(n+1)!} < \frac{3}{(n+1)!}$$

[To see that $e < 3$, notice that $\sum_{k=0}^n \frac{1}{k!} = 1 + 1 + \frac{1}{2} + \frac{1}{2 \cdot 3} + \frac{1}{2 \cdot 3 \cdot 4} + \dots + \frac{1}{n!} <$

$$(1 + 1 + \frac{1}{2} + \frac{1}{2^2} + \frac{1}{2^3} + \dots + \frac{1}{2^n}) - (\frac{1}{2^2} - \frac{1}{2 \cdot 3}) = (3 - \frac{1}{2^n}) - \frac{1}{12}, \text{ and}$$

$$e = \sum_{k=0}^n \frac{1}{k!} + E_n(1) < 3 - \frac{1}{2^n} - \frac{1}{12} + \frac{e}{(n+1)!} < 3 \text{ by choosing}$$

n large enough that $\frac{e}{(n+1)!} < \frac{1}{12}$ — which would happen regardless of how big e was. Alternatively, you could use step functions to show that $\int_1^3 \frac{dx}{x} > 1$.]

This allows us to show that, for example, $\sum_{k=0}^{10} \frac{1}{k!}$ gives at least 6 decimal places of accuracy, since $\frac{3}{11!} \sim 0.00000008$.

② So we can show that certain numbers are irrational.

Theorem 2: Given sequences $a_n, b_n \in \mathbb{Z}$, and $r \in \mathbb{R}$, suppose that for each $\epsilon > 0$ there exists $N \in \mathbb{N}$ s.t.

$$n \geq N \Rightarrow 0 < \left| r - \frac{b_n}{a_n} \right| < \frac{\epsilon}{a_n}.$$

Then $r \notin \mathbb{Q}$.

Proof: If $r = \frac{p}{q}$ ($p, q \in \mathbb{Z}$), then taking $\epsilon = \frac{1}{q}$ gives

$$0 < \left| \frac{p}{q} - \frac{b_n}{a_n} \right| < \frac{1}{q a_n} \Rightarrow 0 < \underbrace{|p a_n - b_n q|}_{\in \mathbb{Z}} < 1 \quad \text{✗} \quad \square$$

$$E_n/r = e, \quad a_n = n!, \quad b_n = \sum_{k=0}^n \frac{n!}{k!}. \quad \text{Put } f(x) = e^x.$$

Of course $b_n/a_n = \sum_{k=0}^n \frac{1}{k!} = T_n(1)$. So the bound on $E_n(1)$ above gives

$$\left| e - \frac{b_n}{a_n} \right| = |E_n(1)| < \frac{3}{(n+1)!} = \frac{3/(n+1)}{n!} = \frac{3/(n+1)}{a_n}.$$

Taking $n \geq N > \frac{3}{\epsilon}$ makes $\frac{3}{n+1} < \epsilon$. Moreover, we know $E_n(1) = \int_0^1 \frac{f^{(n+1)}(t)}{n!} (1-t)^n dt = \frac{1}{n!} \int_0^1 e^t (1-t)^n dt > 0$.

So the hypotheses of Theorem 2 hold, and thus $e \notin \mathbb{Q}$. //

③ In order to compute indeterminate forms (i.e. limits of type $\frac{0}{0}$, 1^∞ , etc.)

We shall make use of the following rules for computing with little- o notation: $o(g(x)) \pm o(g(x)) = o(g(x))$; $o(cg(x)) = o(g(x))$ if $c \neq 0$; $f(x) \cdot o(g(x)) = o(f(x)g(x))$; $o(o(g(x))) = o(g(x))$; and if $\lim_{x \rightarrow a} g(x) = 0$, $\frac{1}{1+g(x)} = 1 - g(x) + g(x) \frac{g(x)}{1+g(x)} = 1 - g(x) + o(g(x))$.

Ex/ Compute $L = \lim_{x \rightarrow 0} \frac{\alpha^x - \beta^x}{x}$. [In the examples, $a = 0$.]

By the Corollary, taking $f(t) = e^t$ we have $E_1(t) = o(t)$.

So $e^t = 1 + t + o(t)$, and substitution yields

$$\alpha^x - \beta^x = e^{x \log \alpha} - e^{x \log \beta} = \cancel{1 + x \log \alpha + o(x \log \alpha)} - \cancel{1 + x \log \beta + o(x \log \beta)} \\ = (\log \frac{\alpha}{\beta}) x + o(x)$$

$$\Rightarrow \frac{\alpha^x - \beta^x}{x} = \log \left(\frac{\alpha}{\beta} \right) + o(1) \xrightarrow[\text{as } x \rightarrow 0]{\text{means } \rightarrow 0} \Rightarrow L = \log \left(\frac{\alpha}{\beta} \right). //$$

Ex / Compute $L = \lim_{x \rightarrow 0} \frac{\log(1+bx)}{x}$ and $L = \lim_{x \rightarrow 0} (1+bx)^{1/x}$, $b \in \mathbb{R}$.

Taking $f(x) = \log(1+bx)$ we have $f'(x) = \frac{b}{1+bx} \Rightarrow (T, f)(x) = bx$.

So $\log(1+bx) = bx + o(x) \Rightarrow \frac{\log(1+bx)}{x} = b + o(1) \Rightarrow L = b$.

Hence also $L = \lim_{x \rightarrow 0} e^{\frac{1}{x} \log(1+bx)} = e^{\lim_{x \rightarrow 0} \frac{1}{x} \log(1+bx)} = e^b$ by

continuity of exp. //

Ex / Find $(T_3 \tan)(x)$.

We have $\sin(x) = x - \frac{x^3}{6} + o(x^4)$, $\cos(x) = 1 - \frac{x^2}{2} + o(x^3)$

$$\Rightarrow \frac{1}{\cos(x)} = \frac{1}{1 - (\frac{x^2}{2} + o(x^3))} = 1 + \frac{x^2}{2} + o(x^2)$$

$$\Rightarrow \tan(x) = \frac{\sin(x)}{\cos(x)} = \left(x - \frac{x^3}{6} + o(x^4)\right) \left(1 + \frac{x^2}{2} + o(x^2)\right) \\ = x + \frac{x^3}{3} + o(x^3). //$$

The Taylor polynomials for $\tan(x)$ more generally are very mysterious.

For instance, T_7 is $x + \frac{2x^3}{3!} + \frac{16x^5}{5!} + \frac{272x^7}{7!}$. The numbers

appearing in the numerators are of course the derivatives of \tan at 0, the even ones being zero. The odd ones are given by a

beautiful formula involving Bernoulli numbers:

$$\tan^{(2m-1)}(0) = (-1)^{m-1} \frac{(4^{2m} - 4^m)}{2m} B_{2m}.$$