

1.(1 pt) The goal of this exercise is to practice finding the inverse modulo  $m$  of some (relatively prime) integer  $n$ . We will find the inverse of 16 modulo 67, i.e., an integer  $c$  such that  $16c \equiv 1 \pmod{67}$ .)

First we perform the Euclidean algorithm on 16 and 67:

$$67 = 4 * \underline{\hspace{2cm}} + \underline{\hspace{2cm}} * 5 + 1$$

[Note your answers on the second row should match the ones on the first row.]

Thus  $\gcd(16,67)=1$ , i.e., 16 and 67 are relatively prime. Now we run the Euclidean algorithm backwards to write  $1 = 67s + 16t$  for suitable integers  $s, t$ .

$$s = \underline{\hspace{2cm}}$$

$$t = \underline{\hspace{2cm}}$$

when we look at the equation  $67s + 16t \equiv 1 \pmod{67}$ , the multiple of 67 becomes zero and so we get  $16t \equiv 1 \pmod{67}$ . Hence the multiplicative inverse of 16 modulo 67 is  $\underline{\hspace{2cm}}$

2.(1 pt) Find the smallest positive integer  $x$  that solves the congruence:

$$10x \equiv 5 \pmod{63}$$

$$x = \underline{\hspace{2cm}}$$

(Hint: From running the Euclidean algorithm forwards and backwards we get  $1 = s(10) + t(63)$ . Find  $s$  and use it to solve the congruence.)

3.(1 pt) We will find the smallest positive integer  $x$  that solves the following system of congruences via the Chinese Remainder Theorem:

$$x \equiv 1 \pmod{3}$$

$$x \equiv 5 \pmod{11}$$

$$x \equiv 12 \pmod{17}$$

$$x \equiv 12 \pmod{47}$$

In the language of Theorem 4 from page 142 of the 4th edition of Rosen, we thus have the values of the following variables:

$$a_1 = 1 \text{ and } m_1 = 3$$

$$a_2 = 5 \text{ and } m_2 = 11$$

$$a_3 = 12 \text{ and } m_3 = 17$$

$$a_4 = 12 \text{ and } m_4 = 47$$

The first step is to calculate  $m = m_1 m_2 m_3 m_4$ . When we do this we get:

$$m = \underline{\hspace{2cm}}$$

The second step is to calculate the  $\hat{m}_k$  which are given by the formula  $\hat{m}_k = \frac{m}{m_k}$ . Enter their values below:

$$\hat{m}_1 = \underline{\hspace{2cm}}$$

$$\hat{m}_2 = \underline{\hspace{2cm}}$$

$$\hat{m}_3 = \underline{\hspace{2cm}}$$

$$\hat{m}_4 = \underline{\hspace{2cm}}$$

Next we find the  $\hat{y}_k$  which are given by solving

$$\hat{y}_k \hat{m}_k \equiv 1 \pmod{m_k}$$

Thus for example, to find  $\hat{y}_1$  we need to solve

$$8789 \hat{y}_1 \equiv 1 \pmod{3}$$

Since we know  $8789 \equiv 2 \pmod{3}$ , this simplifies to

$$2 \hat{y}_1 \equiv 1 \pmod{3}$$

Solve this either by trial and error or by using the Euclidean algorithm and enter the value of  $\hat{y}_1$  below: (Use the canonical representative modulo 3.)

$$\hat{y}_1 = \underline{\hspace{2cm}}$$

Similarly, to find  $\hat{y}_2$  we need to solve

$$2397 \hat{y}_2 \equiv 1 \pmod{11}$$

Since we know  $2397 \equiv 10 \pmod{11}$ , this simplifies to

$$10 \hat{y}_2 \equiv 1 \pmod{11}$$

Solve this either by trial and error or by using the Euclidean algorithm and enter the value of  $\hat{y}_2$  below: (Use the canonical representative modulo 11.)

$$\hat{y}_2 = \underline{\hspace{2cm}}$$

Similarly, to find  $\hat{y}_3$  we need to solve

$$1551 \hat{y}_3 \equiv 1 \pmod{17}$$

Since we know  $1551 \equiv 4 \pmod{17}$ , this simplifies to

$$4 \hat{y}_3 \equiv 1 \pmod{17}$$

Solve this either by trial and error or by using the Euclidean algorithm and enter the value of  $\hat{y}_3$  below: (Use the canonical representative modulo 17.)

$$\hat{y}_3 = \underline{\hspace{2cm}}$$

Similarly, to find  $\hat{y}_4$  we need to solve

$$561 \hat{y}_4 \equiv 1 \pmod{47}$$

Since we know  $561 \equiv 44 \pmod{47}$ , this simplifies to

$$44 \hat{y}_4 \equiv 1 \pmod{47}$$

Solve this either by trial and error or by using the Euclidean algorithm and enter the value of  $\hat{y}_4$  below: (Use the canonical representative modulo 47.)

$$\hat{y}_4 = \underline{\hspace{2cm}}$$

Now that we have all the  $a_k, \hat{m}_k$  and  $\hat{y}_k$ , use the formula  $x = \sum_{k=1}^4 a_k \hat{m}_k \hat{y}_k$  to find an integer solution  $x$  to the original system. The Chinese remainder theorem says that this  $x$  and any integer congruent modulo  $m$  to it, will solve the original system. Enter the SMALLEST positive integer solution to the original system here:  $\underline{\hspace{2cm}}$ .

4.(1 pt) Find the SMALLEST positive integer solution to the following system of congruences:

$$x \equiv 2 \pmod{5}$$

$$x \equiv 1 \pmod{7}$$

The solution is  $\underline{\hspace{2cm}}$ .

5.(1 pt) Use Fermat's Little theorem to compute the following remainders for  $3^{1921}$  (Always use canonical representatives.)

$$3^{1921} = \underline{\hspace{2cm}} \pmod{5}$$

$$3^{1921} = \underline{\hspace{2cm}} \pmod{7}$$

$$3^{1921} = \underline{\hspace{2cm}} \pmod{11}$$

Use your answers above to find the canonical representative of  $3^{1921} \bmod 385$  by using the Chinese Remainder Theorem. [Note  $385 = 5 \cdot 7 \cdot 11$  and that Fermat's Little Theorem cannot be used to directly find  $3^{1921} \bmod 385$  as 385 is not a prime and also since it is larger than the exponent.]  $3^{1921} \bmod 385$  is \_\_\_\_\_

6.(1 pt) Fill in the blanks in the table with the unique integers  $a$  in the range  $0 \leq a \leq 27$  with the given remainders.

Hint: It is probably easiest to just make a table with the numbers between 0 and 27 and their remainders and use that to find the answers. However one can also use the Chinese Remainder Formula

$x = a_1 \hat{m}_1 \hat{y}_1 + a_2 \hat{m}_2 \hat{y}_2$  by finding the  $\hat{m}_k, \hat{y}_k$  once and then plugging in the various remainders for the  $a_k$  to get the various answers.

$a$	$a \bmod 4$	$a \bmod 7$
—	1	3
—	0	2
—	1	1
—	3	2
—	2	5
—	0	2
—	3	6
—	2	0
—	2	3

7.(1 pt) (Modification of exercise 36 in section 2.5 of Rosen.)

The goal of this exercise is to work thru the RSA system in a simple case:

We will use primes  $p = 41, q = 61$  and form  $n = 41 \cdot 61 = 2501$ . [This is typical of the RSA system which chooses two large primes at random generally, and multiplies them to find  $n$ . The public will know  $n$  but  $p$  and  $q$  will be kept private.]

Now we choose our public key  $e = 13$ . This will work since  $\gcd(13, (p-1)(q-1)) = \gcd(13, 2400) = 1$ . [In general as long as we choose an 'e' with  $\gcd(e, (p-1)(q-1)) = 1$ , the system will work.]

Next we encode letters of the alphabet numerically say via the usual:

(A=0, B=1, C=2, D=3, E=4, F=5, G=6, H=7, I=8, J=9, K=10, L=11, M=12, N=13, O=14, P=15, Q=16, R=17, S=18, T=19, U=20, V=21, W=22, X=23, Y=24, Z=25.)

We will practice the RSA encryption on the single integer 15. (which is the numerical representation for the letter "P"). In the language of the book, M=15 is our original message.

The coded integer is formed via  $c = M^e \bmod n$ .

Thus we need to calculate  $15^{13} \bmod 2501$ .

This is not as easy as it seems and you might consider using fast modular multiplication

The canonical representative of  $15^{13} \bmod 2501$  is \_\_\_\_\_